

Optimal adjacency labels for subgraphs of Cartesian products

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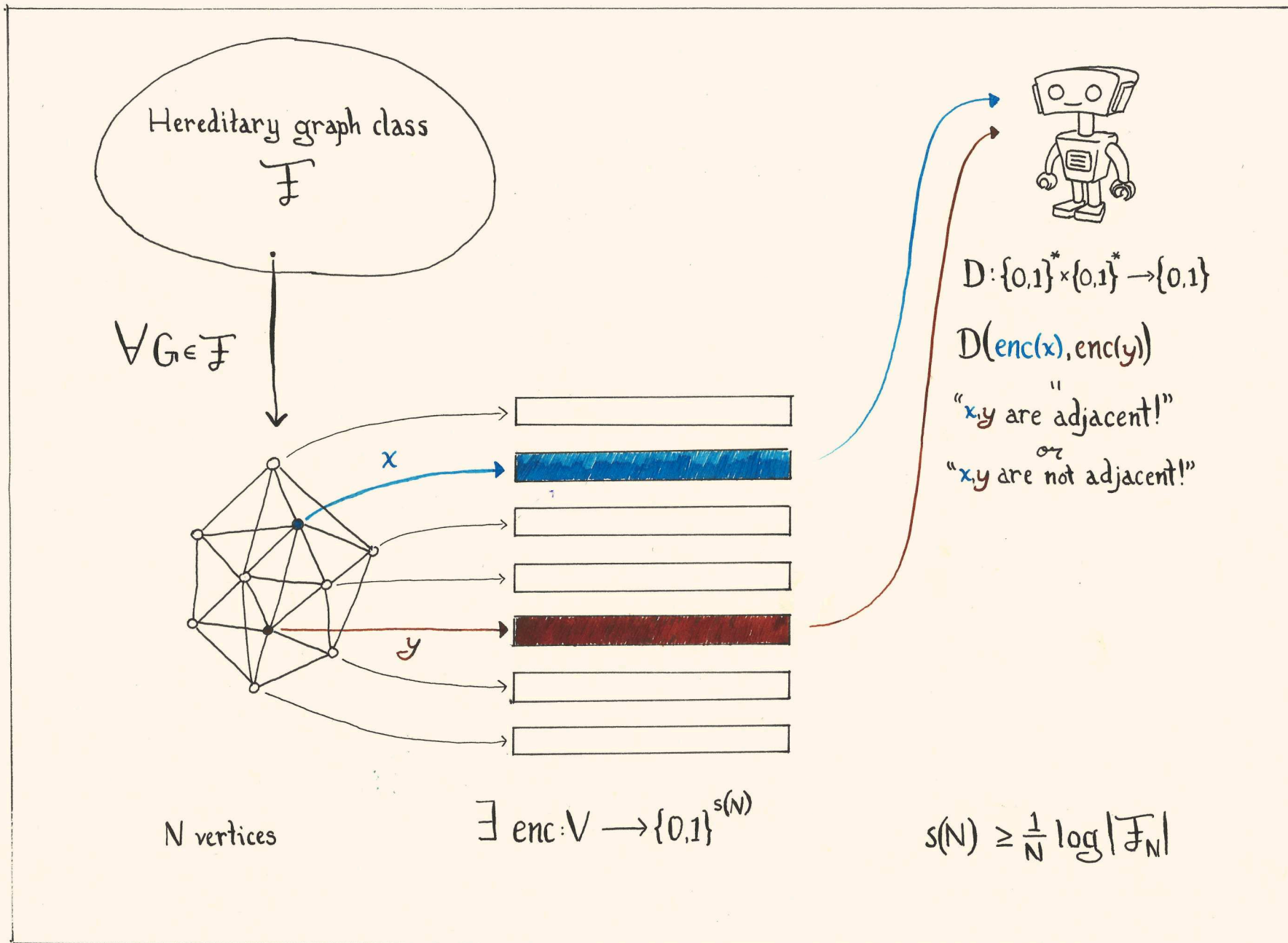
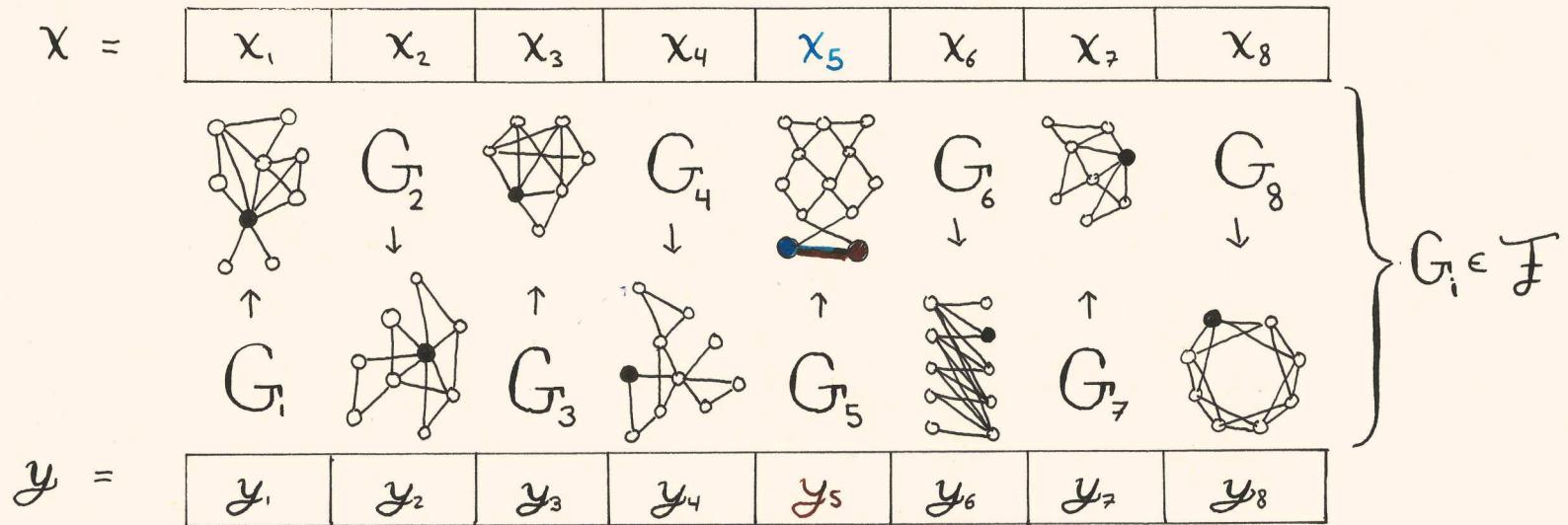


fig 1. ADJACENCY LABELLING

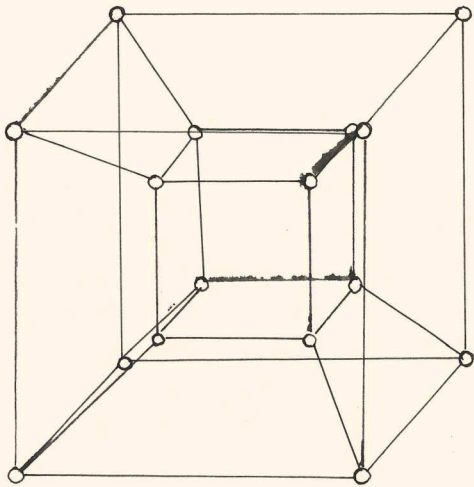
$$\mathcal{F}^\square := \left\{ G = \underbrace{G_1 \square G_2 \square \dots \square G_d}_{\mathcal{F}} \mid d \in \mathbb{N} \right\}$$



$\text{her}(\mathcal{F}^\square)$ | Delete vertices

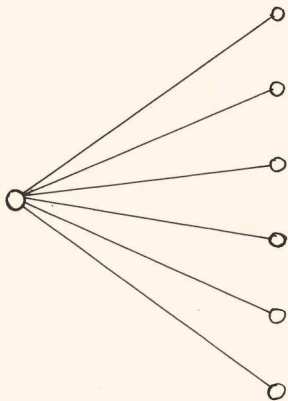
$\text{mon}(\mathcal{F}^\square)$ | Delete vertices & edges

fig2. CARTESIAN PRODUCTS



000000...
 000100...
 001000...
 001100...
 010000...
 010100...
 011000...
 011100...
 ⋮

} $N=2^d$ vertices, $d=\log N$ bits



000000
 000001
 000010
 000100
 001000
 010000
 100000

} $N=d+1$ vertices, $d=N-1$ bits

fig3. Hypercubes

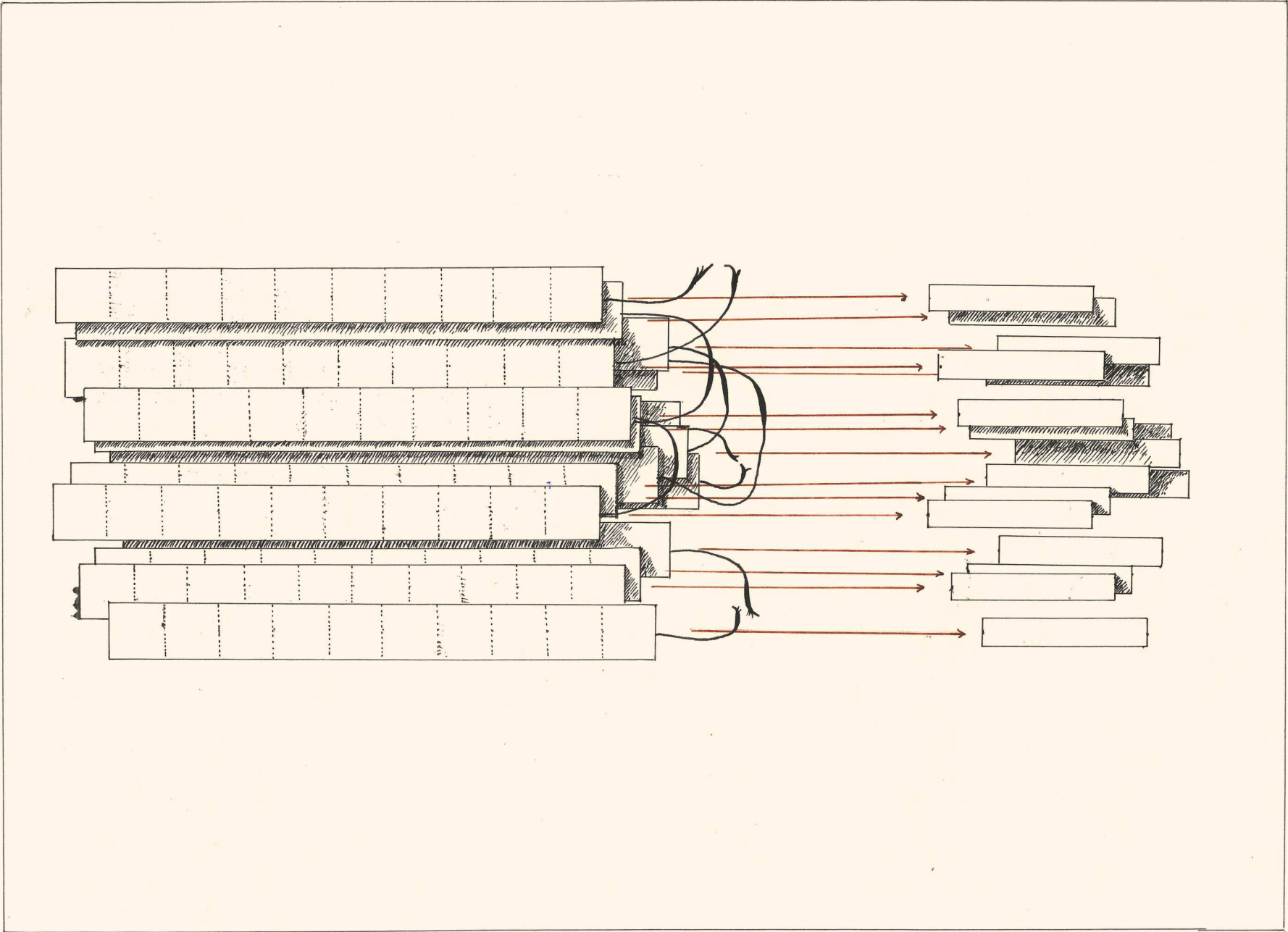
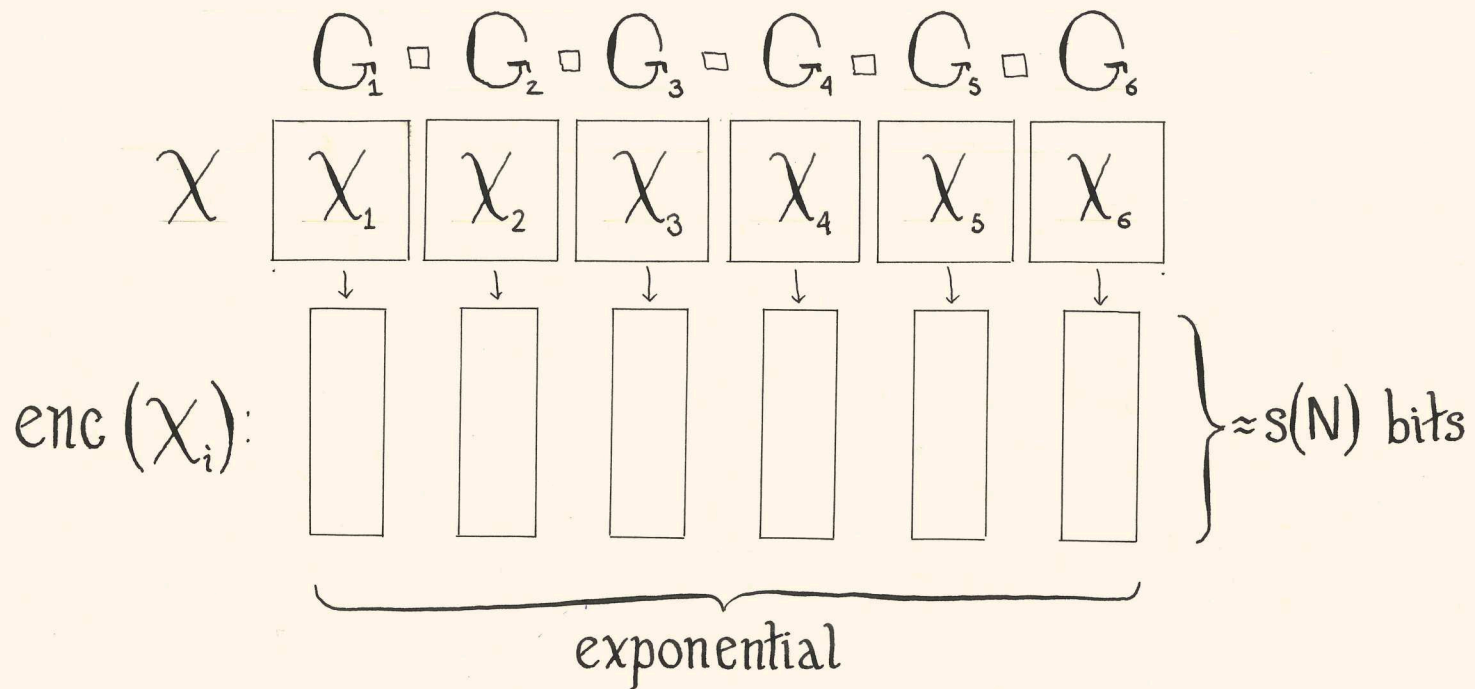


fig 4. ADJACENCY LABELS FOR CARTESIAN PRODUCTS



"Efficient": $s(N) = \Theta\left(\frac{1}{N} \log |\mathcal{F}_N|\right)$

Goal:

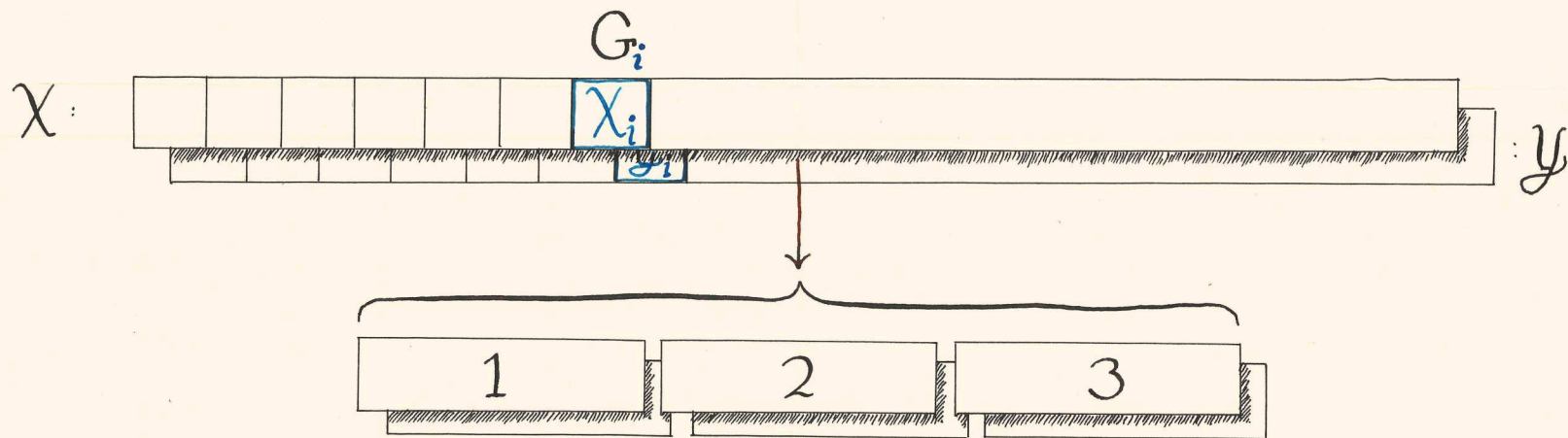
Efficient is impossible for $\text{mon}(\mathcal{F}^\square)$

Why? {
 Chepoi, Labourel, Ratal '20
 Harms, Wild, Zamaraev '22
 Esperet, Harms, Kupavskii '22

Result:

Efficient for $\mathcal{F} \Rightarrow$ Efficient for
 $\text{her}(\mathcal{F}^\square)$ and $\text{mon}(\mathcal{F}^\square)$

fig 5. Results



PHASE 1

x, y differ in only one coordinate, i

Randomized
Communication

PHASE 2

$x_i y_i \in E(G_i)$

Additive
Combinatorics

PHASE 3

$xy \in E(G)$
deleted?

Perfect
Hashing

fig 6. THE PLAN

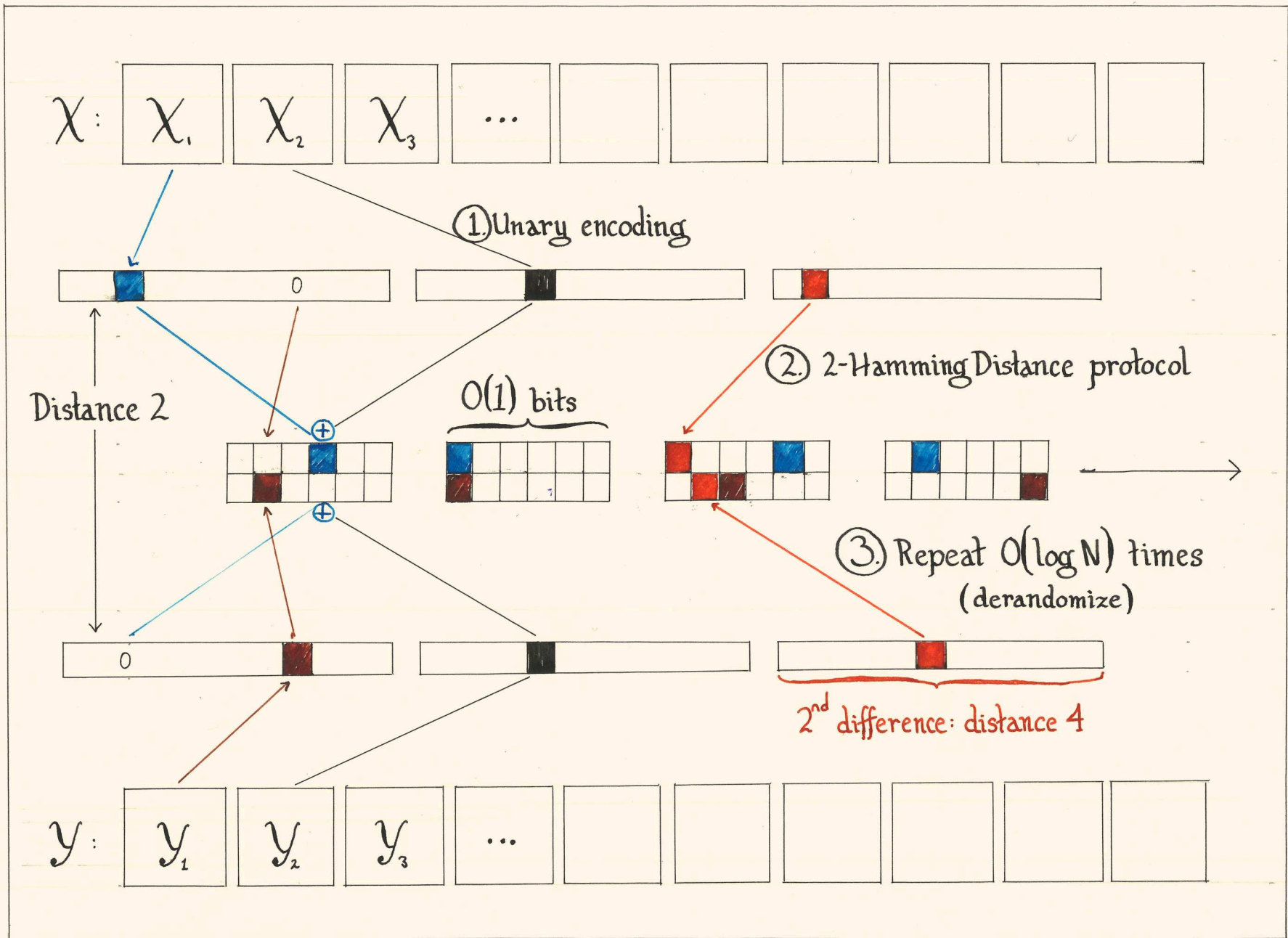


fig 7. PHASE I.

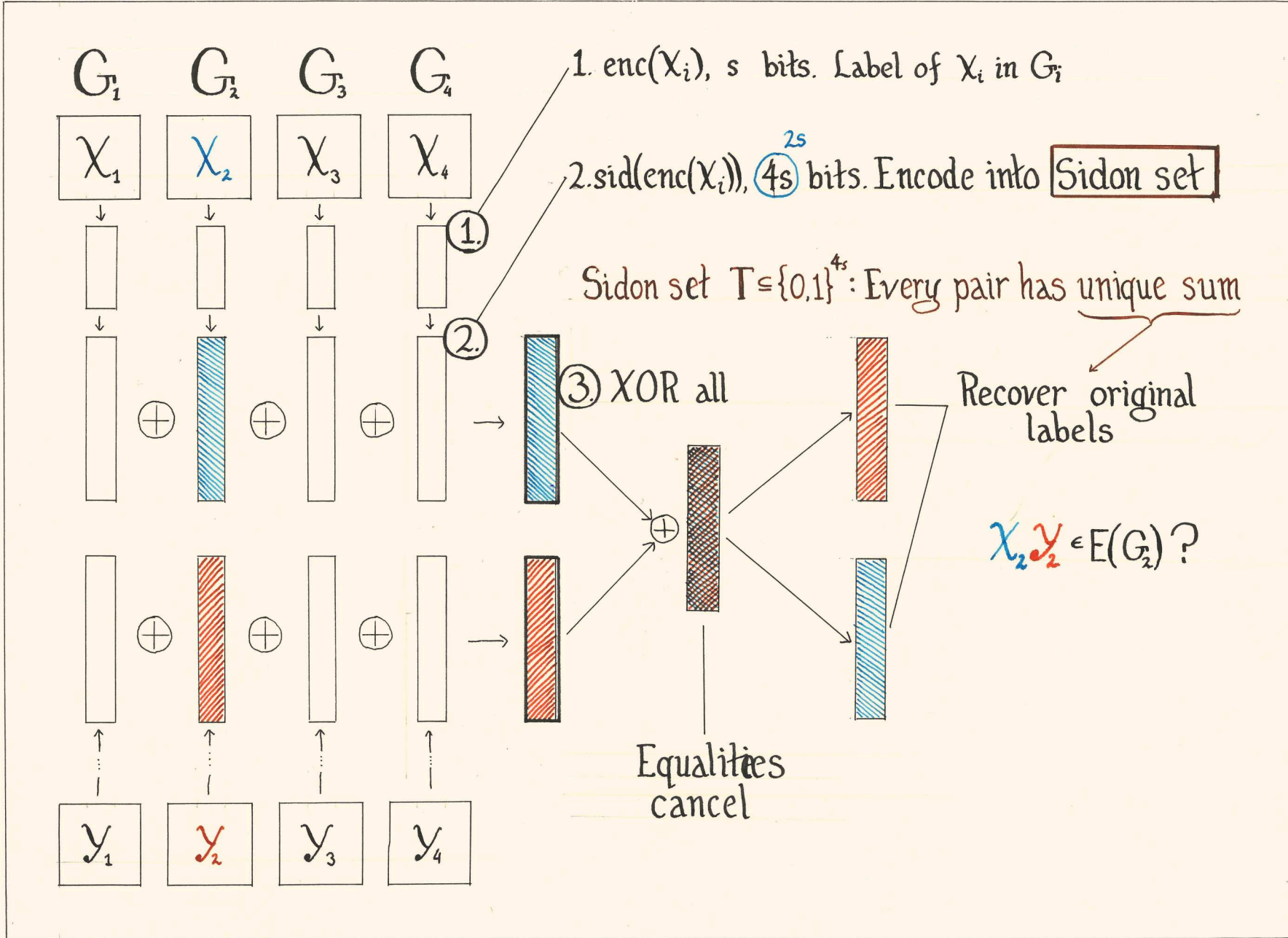
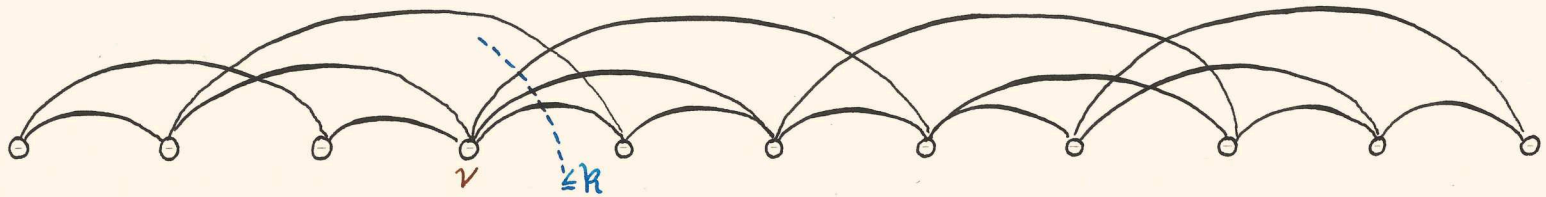
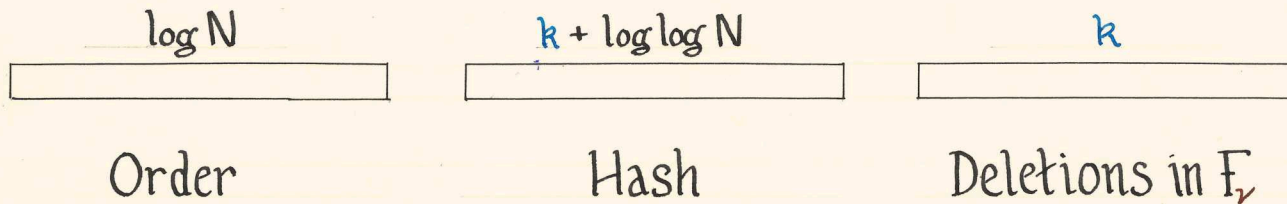


fig 8. PHASE II



Degeneracy k : at most k forward neighbors F_v (Hypercube: $k = \log N$)

Minimal perfect hash: $h_v: F_v \rightarrow [k]$, stored in $O(k + \log \log N)$ bits



$$\underbrace{\log N}_{\text{I.}} + \underbrace{s(N)}_{\text{II.}} + \underbrace{k + \log N}_{\text{III.}}$$

Total size:

(Label size for \mathcal{F})

fig 9. PHASE III.

